

Technologies for Flexible Cross-organizational Case Management Systems

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Overview

- Project Background
- Project Goals
- DCR Graphs
 - Informal introduction
 - Formal semantics
 - Distribution
 - Data and Time
 - Tools
- Future work



Industrial PhD



IT University
of Copenhagen

- Enrolled as PhD Student
- 50% of time (plus 6 months abroad)
- Research, classes, teaching

eXformatics

- Employer (pays salary, travel expenses etc)
- 50% of time
- Tool development, case studies, knowledge dissemination (internal and external presentations, etc)



Danish Government

- Subsidizes project, covers:
- All expenses university
- Approx. half of expenses company



eXformatics

IT University
of Copenhagen

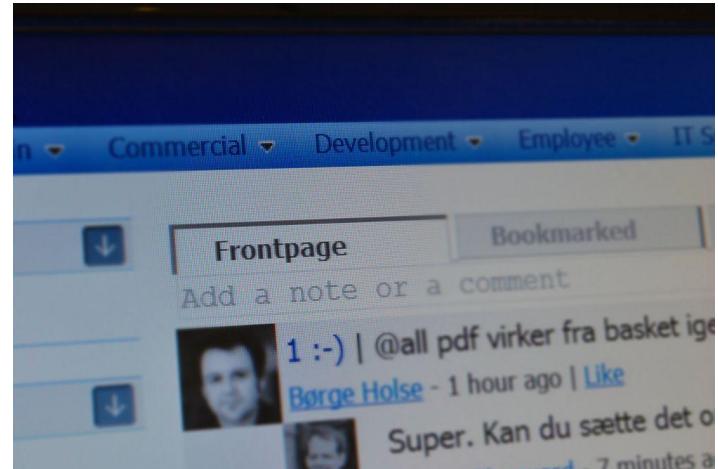
IT University of Copenhagen

- Founded in 1999
- Small independent university (1,500 students)
- Focused on different areas of IT:
 - Communication
 - Business
 - Software Development
 - Games



eXformatics

- Danish company, with a branch in Sweden
- Founded in 2003
- Approx. 20 employees
- Develops Electronic Case- and Document-handling system



Project Goals

Technologies for Flexible Cross-organizational Case Management Systems

- *Case Management* originates from social work and health care, the same concept has since popped up in many other fields under a variety of new names such as *case handling* and *adaptive case management*.



Case Management

- Centered on the concept of a *case*, a process with (usually) the following properties:
 - Aims to solve the needs of a *client* (customer).
 - Goals are *customized* to individual clients and cases.
 - There is a main responsible: the *case worker*.
 - Requires *professional knowledge* and *decision-making*.
 - Involves many participants across multiple divisions and/or organizations



Case Management

- Examples:
 - legal cases in a lawyers firm,
 - claims in an insurance company,
 - hiring of an employee in a HR department,
 - treatment of a patient in a hospital.

Project Goals

Technologies for Flexible Cross-organizational Case Management Systems

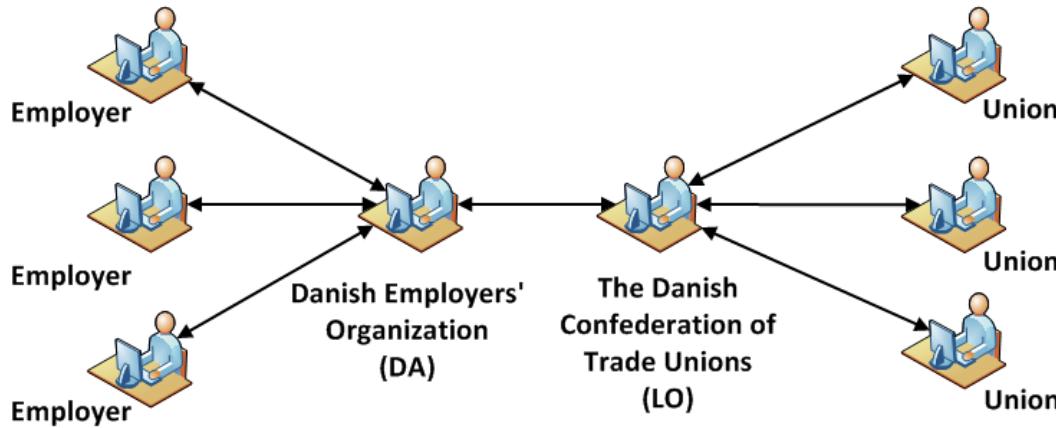
- Caseworkers are *knowledge-workers*: they require flexibility to make their own decisions.
- *Declarative* workflow languages are better at offering this flexibility than *imperative* workflow languages.
- Our declarative model: **DCR Graphs**.



Project Goals

Technologies for Flexible Cross-organizational Case Management Systems

- Case management often involves multiple (sub-) organizations.
- Example: Employer-Employee conflicts in Denmark.



DCR Graphs

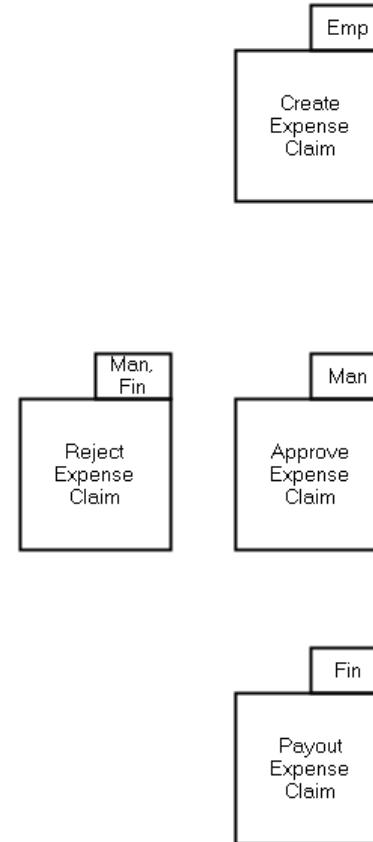
- A declarative workflow language, generalization of *event structures*.
- Consists of *events* and *constraints* between events
- Unconstrained events can happen at any time
- State represented as a *marking* consisting of *executed*, *pending* and *included* events

DCR Graphs

- DCR Graphs are not the first declarative workflow model, but adds:
 - Only 4 basic relations, yet equal to ω -regular languages in expressiveness.
 - Semantics based on transformation of markings: makes execution of DCR Graphs easy to understand and represent graphically.

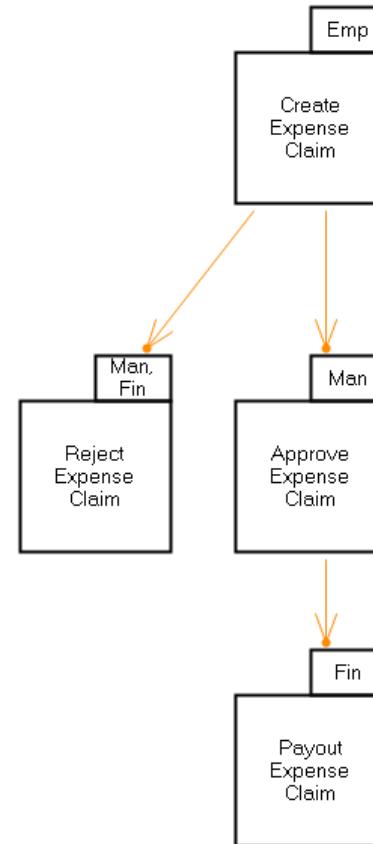
DCR Graphs: Events and Roles

- 4 main activities:
 - Create Expense Claim
 - Approve Expense Claim
 - Reject Expense Claim
 - Payout Expense Claim
- Activities limited to specific roles:
 - Employee
 - Manager
 - Finance Department



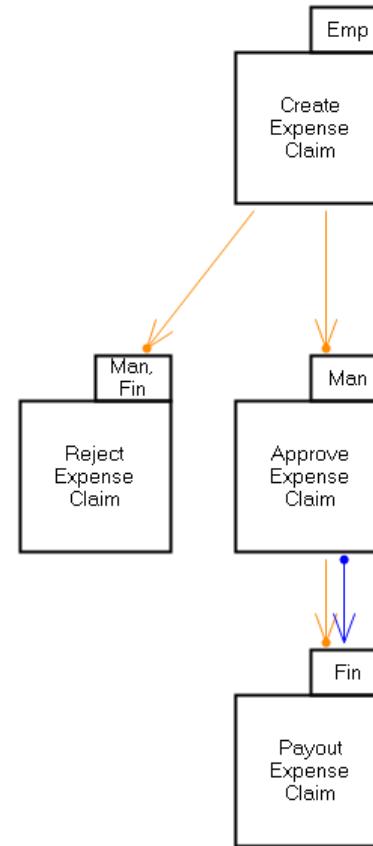
DCR Graphs: Conditions

- A claim should be created before it can be approved or rejected.
- A claim should be approved before it can be paid out.



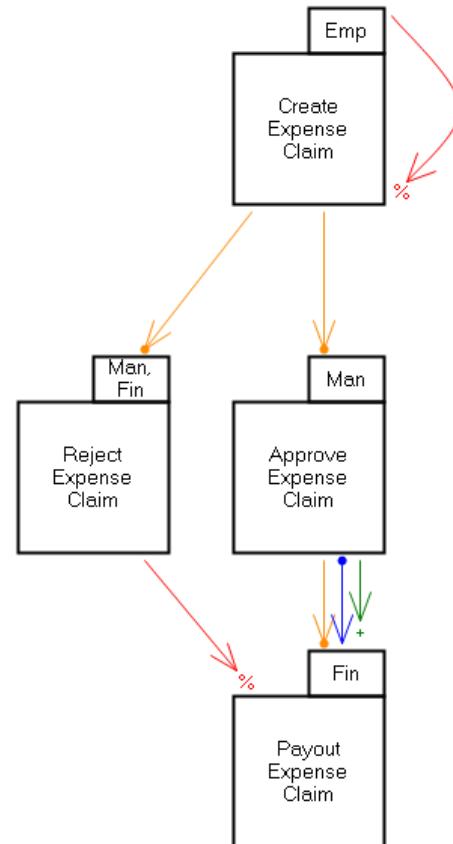
DCR Graphs: Responses

- Once a claim has been approved, it should *eventually* be paid out.



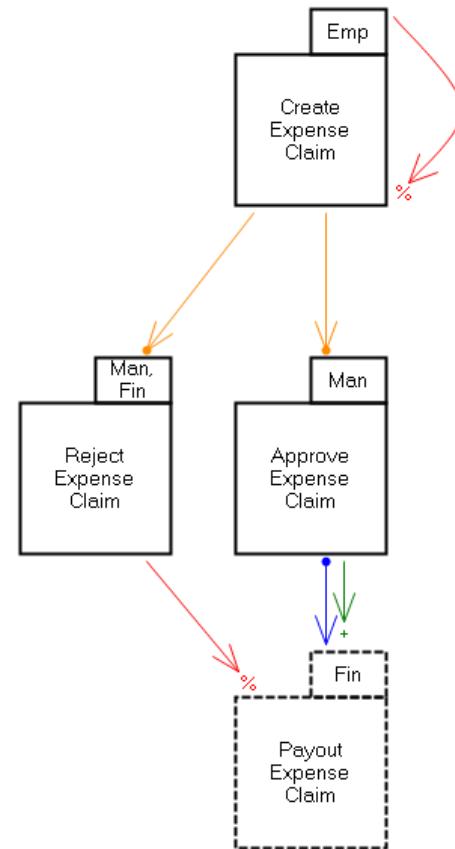
DCR Graphs: Exclusions and Inclusions

- A claim should only be created once. (Every run of the workflow handles a single claim.)
- Once a claim has been rejected, it should not be paid out, unless it is approved again at some later point in time.



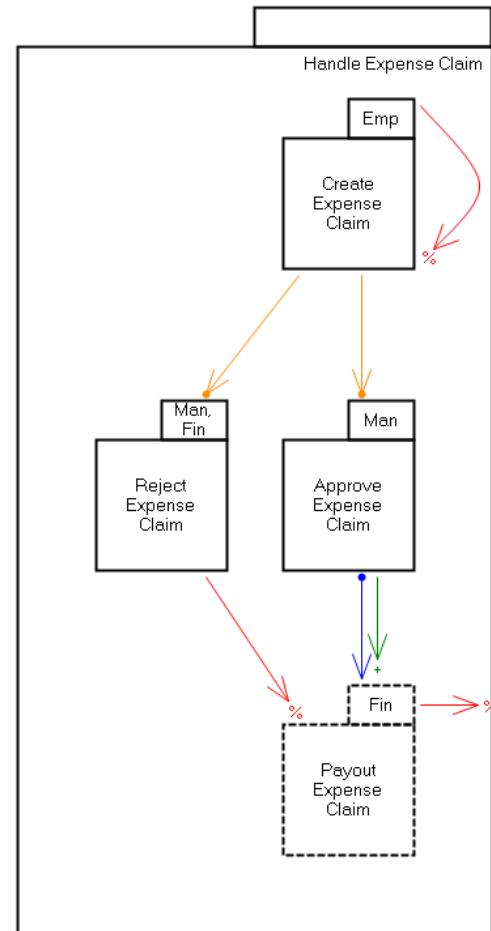
DCR Graphs: Initial Marking

- We can remove the condition from *Approve* to *Payout* if we exclude *Payout* initially.
(*Approve* still needs to be executed before it, because it is the only event that includes it).



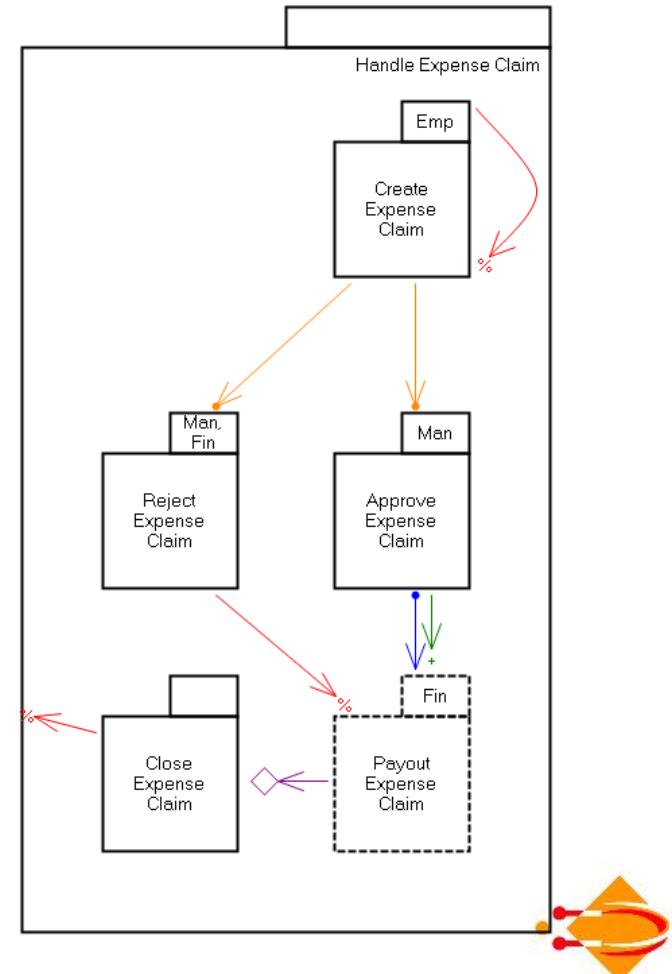
DCR Graphs: Nesting

- Payout should end the process and therefore exclude everything.
- Nesting as a shorthand for having an exclude relation to all four events.



DCR Graphs: Milestone

- Case closed manually, can happen at any time, unless we still need to pay out the claim.
- We will see later that the milestone relation captures the acceptance condition for finite runs.



DCR Graphs: Formal Definition

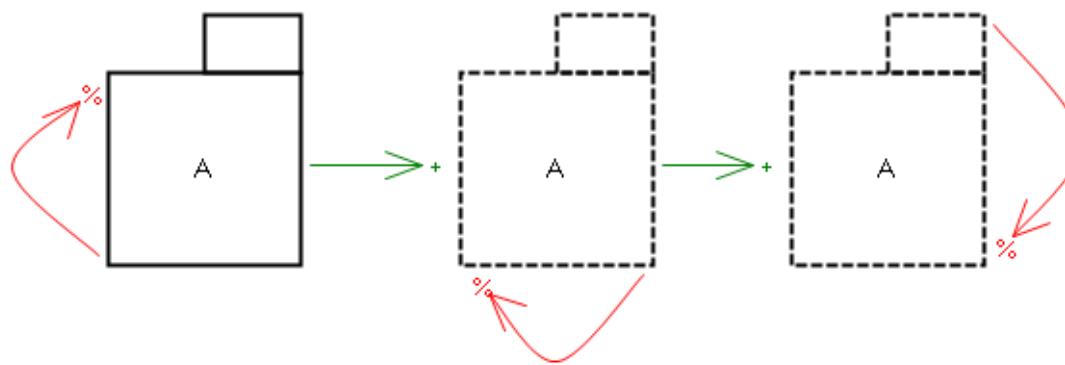
A Nested DCR Graph G is a tuple $G = (\mathcal{E}, \triangleright, M, \rightarrow\bullet, \bullet\rightarrow, \rightarrow\bowtie, \rightarrow+, \rightarrow\%, L, l)$, where

- (i) \mathcal{E} is the set of events
- (ii) $\triangleright : \mathcal{E} \rightharpoonup \mathcal{E}$ is a partial function mapping an event to its super-event
- (iii) $\text{atoms}(\mathcal{E})$ is the set of atomic events that do not have any sub-events
- (iv) $M \in \mathcal{P}(\text{atoms}(\mathcal{E})) \times \mathcal{P}(\text{atoms}(\mathcal{E})) \times \mathcal{P}(\text{atoms}(\mathcal{E}))$ is the initial marking
- (v) $\rightarrow\bullet \subseteq \mathcal{E} \times \mathcal{E}$ is the condition relation
- (vi) $\bullet\rightarrow \subseteq \mathcal{E} \times \mathcal{E}$ is the response relation
- (vii) $\rightarrow\bowtie \subseteq \mathcal{E} \times \mathcal{E}$ is the milestone relation
- (viii) $\rightarrow+ \subseteq \mathcal{E} \times \mathcal{E}$ is the dynamic inclusion relation
- (ix) $\rightarrow\% \subseteq \mathcal{E} \times \mathcal{E}$ is the dynamic exclusion relation
- (x) L is the set of labels (e.g. activities + roles)
- (xi) $l : \mathcal{E} \rightarrow \mathcal{P}(L)$ is a labeling function mapping events to sets of labels.



DCR Graphs: Labels

- Allow multiple events to be mapped to the same label.
- For example allows to express:



- Also introduces non-determinism if only labels are observed (two events with the same label may be enabled at the same time).



DCR Graphs: Formal Semantics

- Nested DCR Graph is flattened into a regular DCR Graph (relations go from/to all sub-events and super-events are discarded).
- An event is *enabled* if:

$$M \vdash_G e, \text{ if } e \in In \text{ and } (In \cap \rightarrow^\bullet e) \subseteq Ex \text{ and } (In \cap \rightarrow^\diamond e) \subseteq E \setminus Re$$

↑
Event is included

↑
All *included* condition events
have been executed

↑
None of the included
milestone events is a pending
response

DCR Graphs: Formal Semantics

- Result of executing an event:

$$(Ex, Re, In) \oplus_G e = (Ex \cup \{e\}, (Re \setminus \{e\}) \cup e \xrightarrow{\bullet}, (In \setminus e \rightarrow \%) \cup e \rightarrow +)$$

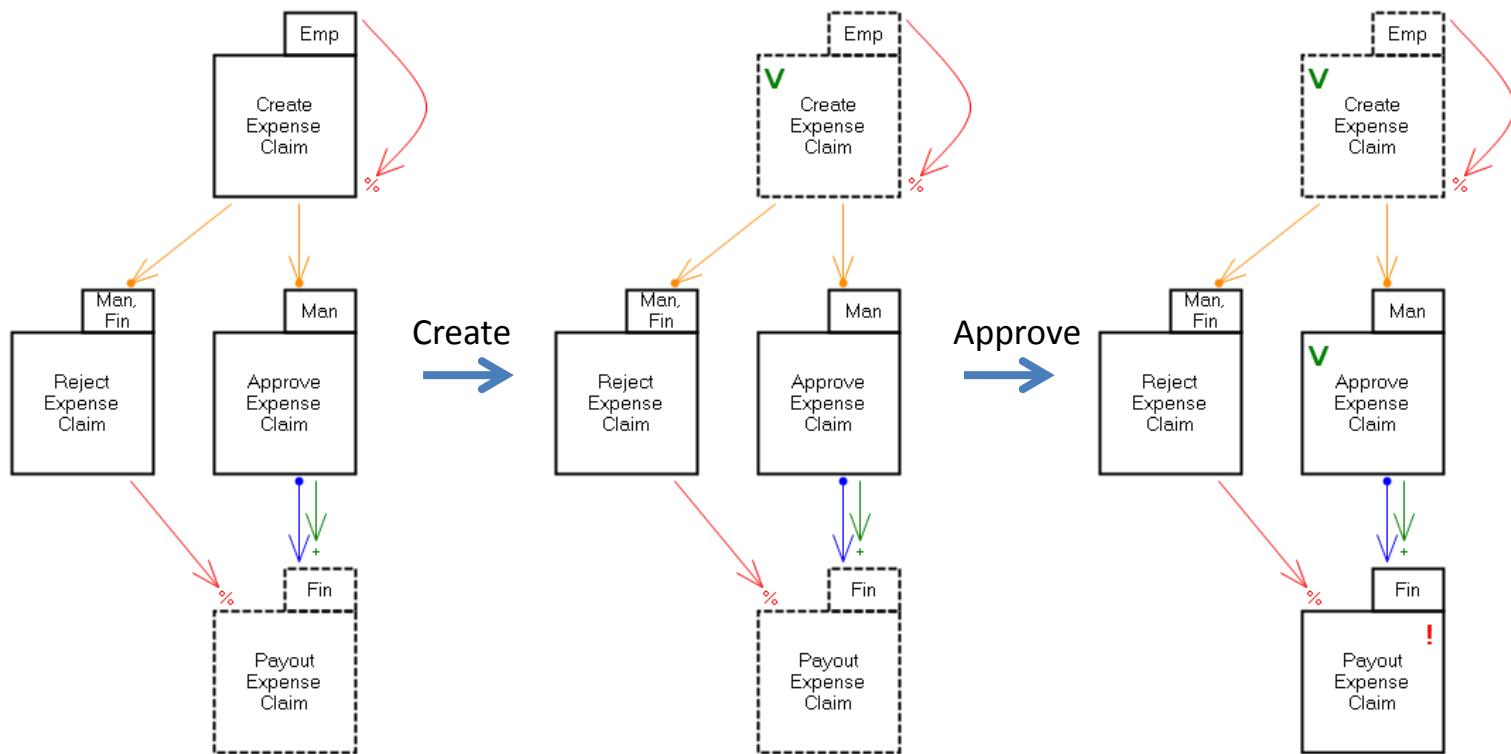
Event e is added to the set of executed events

Event e is removed from the set of pending responses. Afterwards all events that are a response to e are added to the set of pending responses.

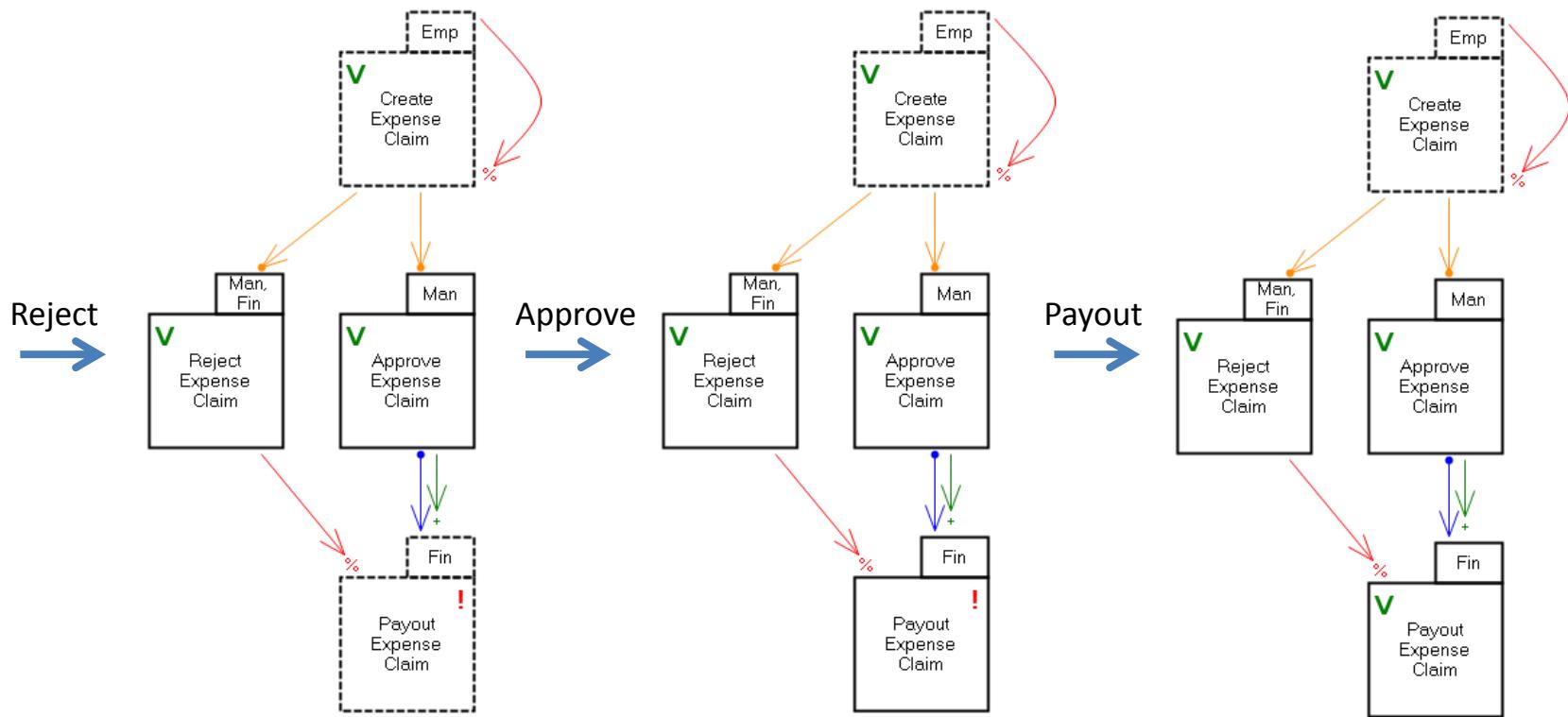
All events that are excluded by e are removed from the set of included events. Afterwards all events that are included by e are added to the set of included events.

- Note that:
 - If an event is a response to itself, it will remain in the set of pending responses after execution.
 - If an event e both includes and excludes an event e', e' will always be included after execution of e.

DCR Graphs: Example Run

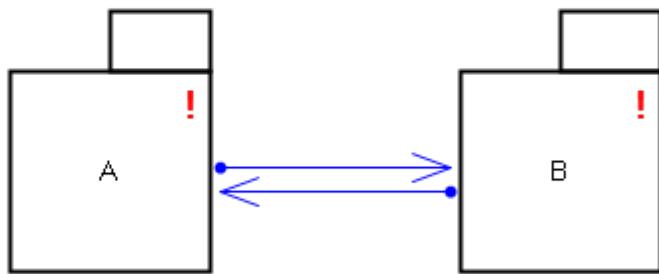


DCR Graphs: Example Run



DCR Graphs: Accepting runs

- A run is accepting if every pending event is eventually executed or excluded.
- Finite runs: Any marking where none of the included events are also pending is accepting.
- Infinite runs:

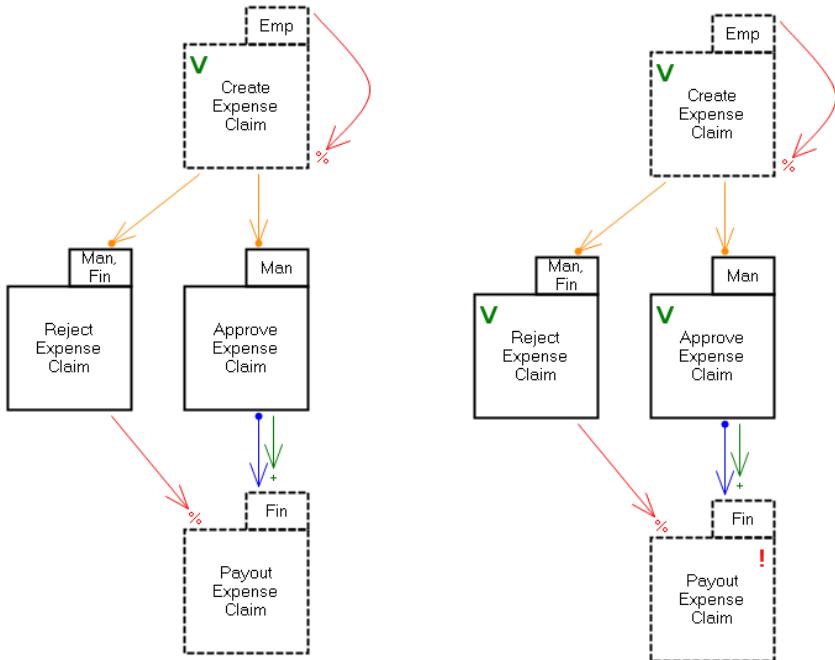


A^ω is not accepting, but $(AB)^\omega$ is.

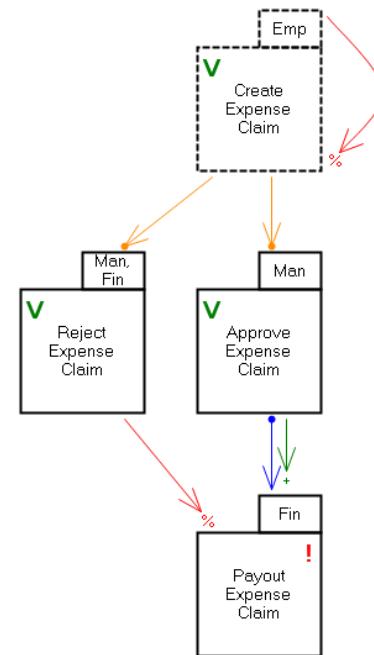


DCR Graphs: Accepting markings

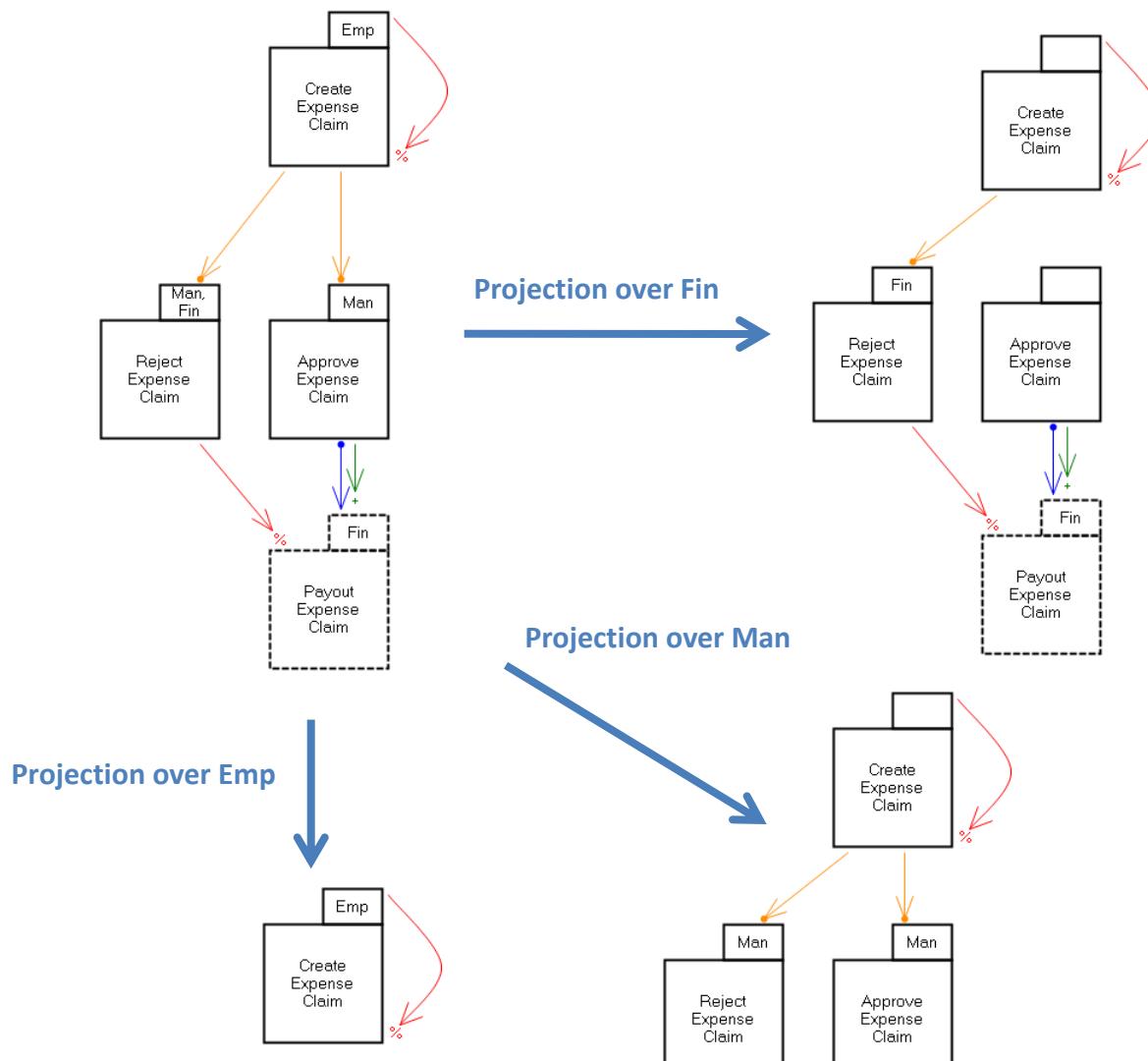
Accepting:



Not Accepting:

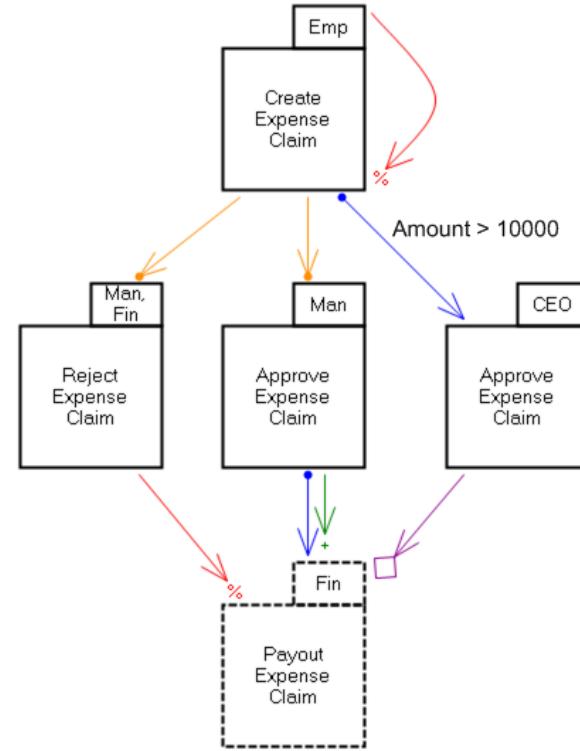


DCR Graphs: Distribution



DCR Graphs: Global Data

- Global data store
- Events and relations can be guarded by expressions on variables in the data store.
- Example: If the amount of an expense claim is larger than 10000 euros, the CEO needs to approve it as well.

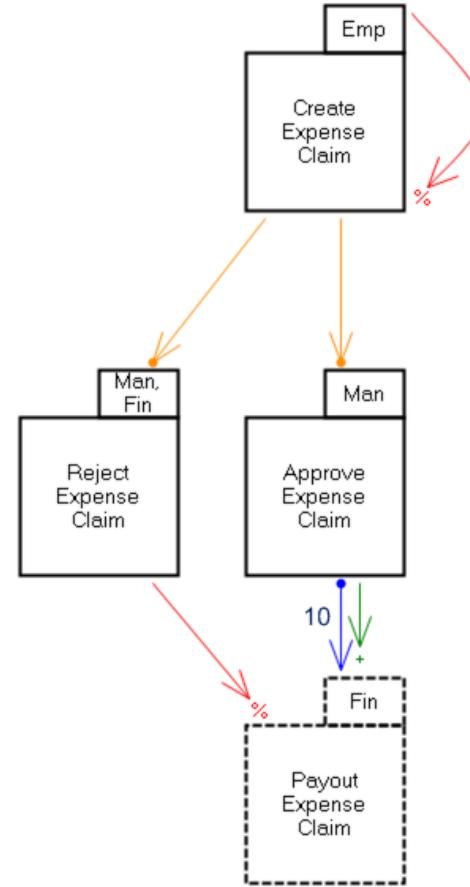


DCR Graphs: Parameterized Events

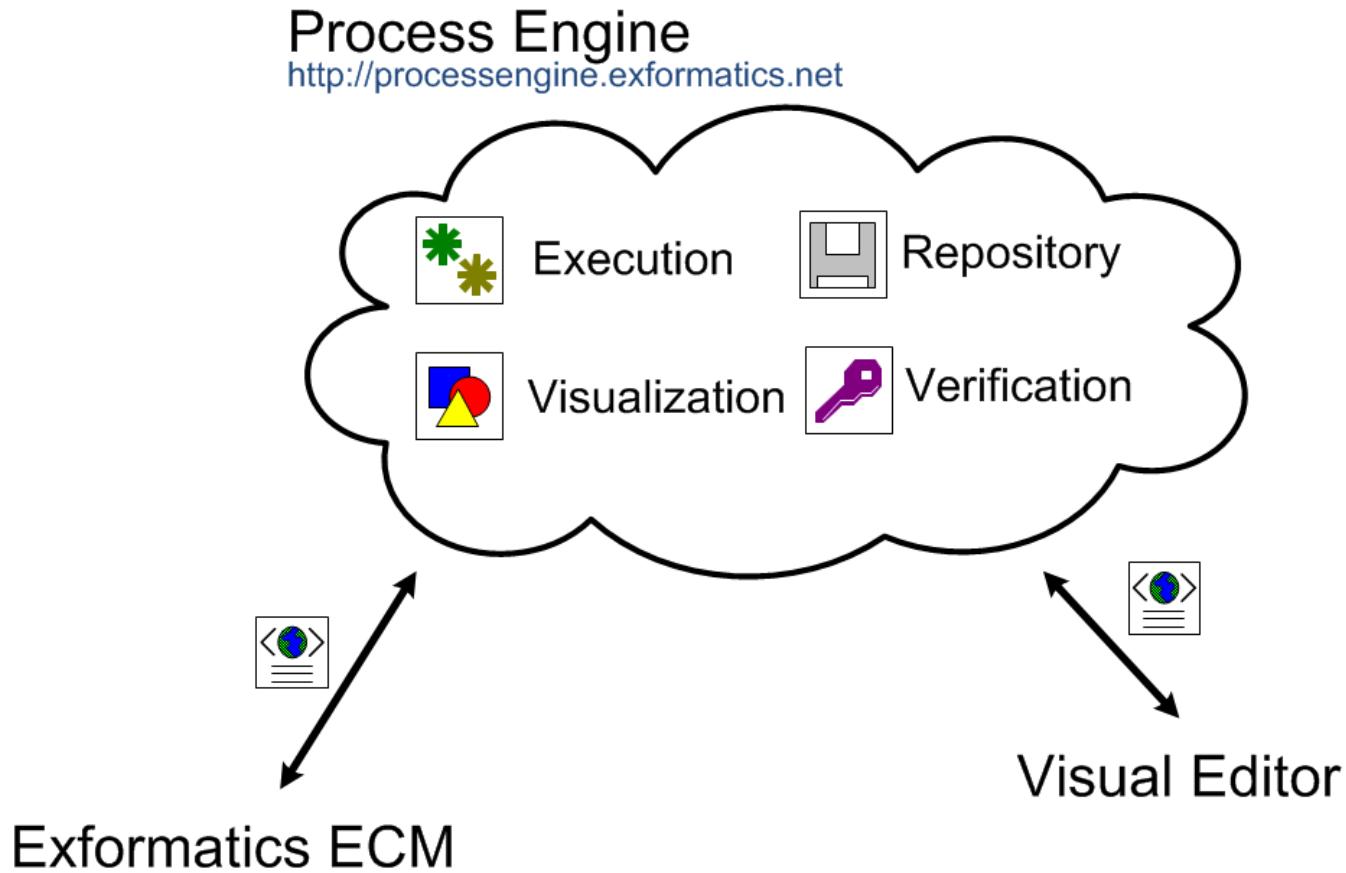
- Events are parameterized, for example: *Create Expense Claim (string Name, int Amount)*
- Relations can be parameterized, for example the *Create Expense Claim* event above can be a condition for an *Approve Expense Claim* event with the same *Name* and *Amount*.

DCR Graphs: Time

- A discrete global clock
- Execution step: an event, or a time step.
- Conditions and responses guarded by time-bounds.
 - Condition: Minimum time that should have passed since the last time an event happened.
 - Response: Deadline for when something should happen.
- Example: After approval, an expense claim should be paid out within 10 time steps.



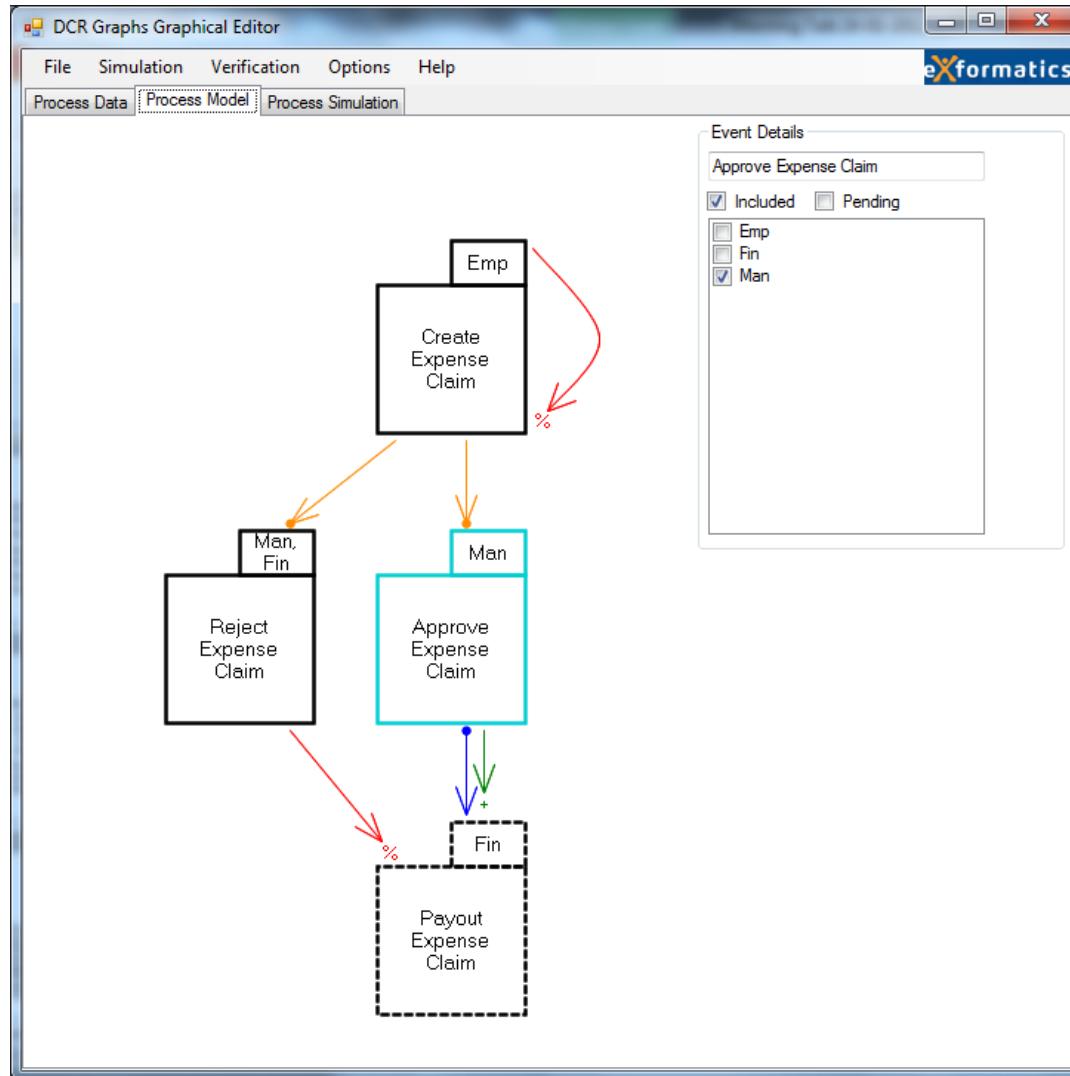
DCR Graphs: Tools



Exformatics ECM

Visual Editor

DCR Graphs: Tools



Future Work

- Extend distributed semantics (asynchronous execution).
- Types for DCR Graphs, based on session types
- Disjunctive DCR Graphs (require one of two options as a response)
- Extend tools to handle some of the concepts discussed before

Future Work

- Explore mappings to other formalisms:
 - Declare: A declarative workflow language [Aalst et al.]
 - Guard-Stage-Milestone model: data-centric workflow model with declarative elements [Hull et al.]
 - Petri nets
- Overall goal: makes DCR Graphs more interesting commercially as it allows for interoperability with other formalisms and tools. (In particular BPMN.)

Conclusions

DCR Graphs:

- Few simple primitives, yet expressive
- Combines declarative specification with effective execution
- Support for nesting and distribution.
- Support for data and time.
- Broad application domain (healthcare workflows, case management, complex event processing)
- Positive feedback from industry
- Thank you for your time! Questions?

